

Warm-up exercise

- A riddle courtesy of Dr. Chris Johnson: what will the following Java program print?
 - (and why?)

```
public class Riddle
{
    public static void main(String[] args) {
        System.out.println(getTrue() || getFalse());
    }
    public static boolean getTrue() {
        System.out.print("true ");
        return true;
    }
    public static boolean getFalse() {
        System.out.print("false ");
        return false;
    }
}
```

Warm-up exercise

- A riddle courtesy of @cpuGoogle: what will the following C++ program return?
 - (assume “=” is evaluated R-to-L and returns lvalue)

```
int main() {  
    int i = 0;  
    (i = 12) = i++;  
    return i;  
}
```

CS 430 Spring 2022

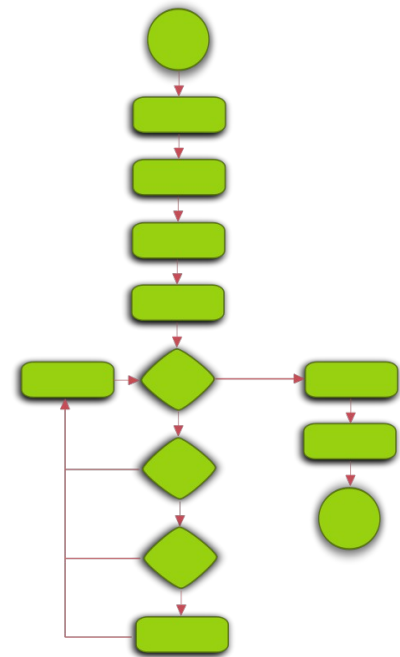
Mike Lam, Professor

$$x = a + b + c$$

$$y = \sin(x)$$

$$E = mc^2$$

$$e^{ix} = \cos x + i \sin x$$



Expressions and Control Structures

Motivation

- Most programs produce and consume information
 - Code that produces information is called an **expression**
 - Evaluating an expression produces a **value**
- Most programs require decision-making and repetition
 - Dynamic **control flow** enabled by **conditional jumps** at the machine level (remember CS 261?)
 - In high-level languages, control flow is prescribed by **control structures**

Expressions

- **Expression**: specification of computation
 - Fundamental to high-level languages
 - Form/syntax expressed using BNF grammars
 - Four main components:
 - 1) Operations
 - 2) Operands
 - 3) Parentheses
 - 4) Function calls

Expressions

- **Operators**: symbols representing computation
 - Number of operands: **unary** vs. **binary** vs. **ternary**
 - Unary examples: `!x` `x++`
 - Binary examples: `x+y` `x=y` `x+=y`
 - Ternary example: `x?y:z`
 - Code location relative to operands: **infix** vs. **prefix** vs. **postfix**
 - Infix examples: `x*y` `x `mod` y`
 - Prefix examples: `~x` `*x` `++x`
 - Postfix examples: `x?` `x++`
 - **Precedence** (higher or lower)
 - **Associativity** (left or right)

Expressions

- **Overloaded** operators: similar operators on different data types
 - Most languages provide overloaded operators for built-in types
 - E.g., '/' is used for both integer and floating-point division
 - Some languages allow overloaded operators for user-defined types
 - **Type parity**: user-defined types are indistinguishable from built-in types

```
class Vector2D {
public:
    double x, y;

    Vector2D(double x, double y) {
        this->x = x;
        this->y = y;
    }

    Vector2D operator+(const Vector2D& v) {
        return Vector2D(this->x + v.x,
                        this->y + v.y);
    }

    Vector2D operator*(int scale) {
        return Vector2D(this->x * scale,
                        this->y * scale);
    }
};
```

```
#include <iostream>
using namespace std;

ostream& operator<<(ostream& os, const Vector2D& v) {
    os << "(" << v.x << "," << v.y << ")";
    return os;
}

int main() {
    Vector2D v1(1.0, 2.0), v2(2.8, 3.5);
    Vector2D v3 = v1 + v2;
    Vector2D v4 = v1 * 3;
    cout << v3 << endl << v4 << endl;
    return 0;
}
```

```
Output:
(3.8,5.5)
(3,6)
```

Expressions

- **Short-circuited** operators (e.g., `&&`, `||`, `?:`)
 - At least one operand is not evaluated if not needed
 - **Optional chaining** operators (not covered in CPL) guard access to potentially-null pointers/references
 - **Null coalescing** operators (also not in CPL) allow for default values
 - We'll see another (more functional) take on this later

```
if (pet !== null) {  
  console.log(pet.name);      // what if "pet" is null?  
}
```

```
// optional chaining operator '?'  
console.log(pet?.name);
```

```
// null coalescing operator  
console.log(pet?.name ?? 'NPE: no pet exception');
```


Expressions

- **Operands**: input data for computation
 - **Evaluation order** (left-to-right or right-to-left)
 - Errors
 - Overflow and underflow
 - Division by zero
 - Floating-point issues (e.g., NaN, subnormal)
 - Type conversions
 - **Narrowing** vs. **widening**
 - **Implicit** vs. **explicit**

Type Compatibility

- **Type conversions**
 - If variables of a particular type can be converted to a different type, those types are **compatible**
 - **Widening** vs. **narrowing**
 - Latter may cause information loss
 - (in CPL definition, former may cause *precision* loss but not *accuracy* loss; e.g., int to float)
 - **Implicit** vs. **explicit**
 - Implicit: **coercion**, e.g., `float x = 5;`
 - Explicit: **casting**, e.g., `int x = (int)3.14;`

Expressions

- Parentheses
 - Explicit precedence and associativity
 - **Tuple** creation
 - Function invocation (in some languages)
- Function calls
 - **Side effects**: a function changes a parameter or a non-local variable
 - **Referential transparency**: expressions with the same value can be substituted for each other (i.e., no side effects)

`(2 + 2)`

`(2 * 2)`

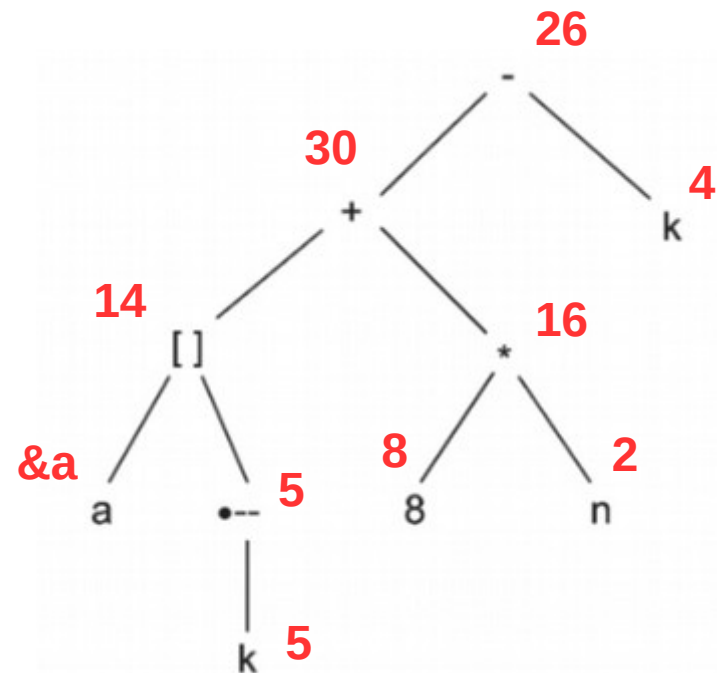
```
def foo
  $n += 1
  return 4
```

Assignment Statements

- Symbol and ambiguity with equality operator
 - "=" vs. "!=" vs. "==" vs. "←"
 - Assignments as expressions; returns lvalue or rvalue?
- **Conditional targets** (ternary LHS)
 - $(n > 5 ? a : b) = n * 2$
- **Compound assignments**
 - Shortened forms of an assignment: "+=" and "++"
 - Example of **syntactic sugar** (unnecessary but helpful)
- **Multiple assignments**
 - $a, b = c/2, c\%2$ $a, b = b, a$

Evaluating Expressions

- Construct expression tree (e.g., *parse* the expression)
 - Build tree from root (“which operation is done last?”)
 - Use precedence and associativity to guide you
- Evaluate using a post-order traversal
 - Use evaluation order to guide you
 - Track side effects as you go
- Example:
 - Use standard prec., assoc. and eval order
 - Suppose $k=5$, $n=2$
 - Suppose $a=\{1, 3, 6, 8, 11, 14, 16\}$
 - Evaluate $a[k--]+8*n-k$



$k = \cancel{4} 5$

Control Structures

- **Control flow path**: sequence order of executed instructions
- **Control structure**: control statement and its associated flow path
- **Selection** statements (e.g., `if/then/else`, `switch/case`)
 - Choose between alternative control flow paths
- **Iteration** statements (e.g., `do`, `while`, `for`, `until`)
 - Repeatedly execute a control flow path
- How many kinds of control statements?
 - Many: higher expressivity
 - Few: higher readability, learnability, and orthogonality

Selection Structures

- **Two-way selection** (`if/then`)
 - Inclusion of "else" clause
 - Blocks often delimited by braces, keywords (e.g., "begin", "end") or indentation
 - Nesting issues
- **Multiple selection** (`switch/case`)
 - Form ("`if/elseif/else`" vs. "`switch/case`")
 - Case value types
 - Multiple execution
 - Fallthrough
 - Default cases
 - Efficient implementation using jump tables

Iteration Structures

- Control form: **logic** vs. **counter** vs. **user-controlled** vs. **iterator-based**
 - Counter loop parameters: loop variables, initial/terminal values, step sizes
 - Counter variable in scope outside loop? (no, starting with Ada)
- Control location: **pre-test** vs. **post-test** vs. **user-defined**
- Examples:
 - While loop: **logic pre-test**
 - Do-while loop: **logic post-test**
 - For loop: **counter pre-test**
 - For-each “enhanced for” loop: **iterator-based pre-test**
- Functional languages: recursion instead of iteration

Language Design

- Can iteration structures have multiple entries?
 - General answer: no!
 - Increase in flexibility/expressiveness is small relative to decrease in readability
- Can iteration structures have multiple exits?
 - For most procedural languages: yes
 - Same as "should goto or break be included?"

Minimally-Sufficiency Constructs

- Böhm and Jacopini (1966)
 - “Structured program theorem”
 - Strictly necessary: 1) sequencing, 2) two-way logical selection, and 3) logical iteration
 - Can implement ALL flowchart-representable programs
 - Alternatively: a selectable goto statement
 - E.g., “if (E) goto L1” **goto code** or **conditional branches** from CS 261!
 - Could be a backwards jump for iteration
 - Facilitates automated translation of block-structured code
 - Conversion to **linear** code (e.g., machine code)
 - Use “templates” to guide translation

Minimally-Sufficient Constructs

if statement: if (E) B1

```
<< E code >>
```

```
if E goto l1
```

```
goto l2
```

```
l1:
```

```
<< B1 code >>
```

```
l2:
```

Minimally-Sufficient Constructs

if statement: if (E) B1 else B2

```
<< E code >>
```

```
if E goto l1
```

```
goto l2
```

```
l1:
```

```
<< B1 code >>
```

```
goto l3
```

```
l2:
```

```
<< B2 code >>
```

```
l3:
```

Minimally-Sufficient Constructs

while loop: while (E) B

```
l1:                                ; CONTINUE target
    << E code >>
    if E goto l2
    goto l3
l2:
    << B code >>
    goto l1
l3:                                ; BREAK target
```

Minimally-Sufficient Constructs

for loop: for V in E1, E2 B

<< E1 code >>

<< E2 code >>

V = E1

l1:

if (V >= E2) goto l2

<< B code >>

V = V + 1

; CONTINUE target

goto l1

l2:

; BREAK target

Minimally-Sufficient Constructs

- Use only the following constructs:

- $S \rightarrow S; S$
- $S \rightarrow \mathbf{if} (E) \{ S \} \mathbf{else} \{ S \}$
- $S \rightarrow \mathbf{while} (E) \{ S \}$
- $S \rightarrow \langle \textit{assignment statement} \rangle$
- $E \rightarrow \langle \textit{boolean expression} \rangle$

- Rewrite the following Ruby code:

```
until a >= b
  a += 5
end
```

```
3.times do
  x = x * 2
end
```

```
1.upto(10) do |i|
  y = y + i
end
```

```
case (n % 3)
when 0
  d = 1
when 1
  d = 2
when 2
  d = 3
end
```

```
if x > 90 then
  g = 'A'
elsif x > 80 then
  g = 'B'
elsif x > 70 then
  g = 'C'
else
  g = 'D'
end
```

Greatest Argument in PL History

- "Should languages provide a goto statement?"
 - Pro: extremely powerful construct – high expressiveness and writability
 - Against: without restrictions, can make programs very difficult to understand – low readability and maintainability
- Classic 1968 CACM letter by Edsger Dijkstra: "*Go To Statement Considered Harmful*"
 - Widely misunderstood
 - Original title: "A Case Against the Goto Statement"
 - Criticized **excessive** use of goto
 - Consensus: structured control flow is safer
 - Use control structures, exceptions, or tail recursion instead
 - Only C descendants tend to have goto statements these days

Guarded Commands

- Dijkstra (1975): **guarded** selection and iteration statements: `if/fi` and `do/od`
 - More than one boolean condition may be true
 - Control flow path is chosen **non-deterministically** out of the available true conditions
 - Pro: sometimes more elegant and easily proven correct
 - Con: greatly-increased complexity and lowered readability

Guarded Commands

- Maximum of (x,y):

```
if  x >= y  →  max := x
[]  y >= x  →  max := y
fi
```

- Sorting four integers (q1, q2, q3, q4):

```
do  q1 > q2  →  temp := q1; q1 = q2; q2 := temp;
[]  q2 > q3  →  temp := q2; q2 = q3; q3 := temp;
[]  q3 > q4  →  temp := q3; q3 = q4; q4 := temp;
od
```